

# Characterisation of qutrit universal gates

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- How to estimate the performance of quantum gates?
- Why is randomised benchmarking chosen?
- Universal qutrit randomised benchmarking.

Qutrits = three level systems =  $\text{span}(|0\rangle, |1\rangle, |2\rangle)$  **single particle**

- Larger Hilbert space.
- More natural: avoid the truncation of a native higher level system.
- Multiple implementations: (superconductors, ion traps, and more) [3, 6, 13, 7, 4, 14, 5, 10, 2, 9].

## Gates

- 1 Group element  $g \in G$ .
- 2 Unitary representation  $U_g$  labelled by a group element.
- 3 Unitary (Schroedinger) evolution:  $\rho \mapsto U_g \rho U_g^\dagger$ .
- 4 Physical implementation labelled by a group element:  $\boxed{g}$ .
- 5 Mathematical modelling by a quantum channel:  $\tilde{U} := \mathcal{E}(U_g \rho U_g^\dagger)$ .

# Estimating performance = characterising

## Why?

- Estimate the performance to improve the gates.
- Know how many operations we can perform before the computation is purely random.

## Why is characterising not easy?

- Stochastic nature of quantum operations.
- Tomography (reconstruction) is expensive (wrt scaling).

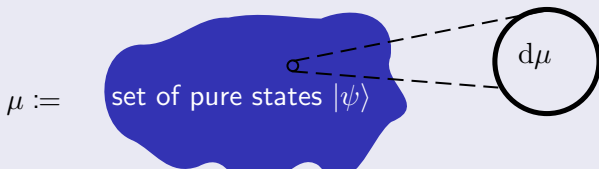
# Fidelity

## Fidelity between states

- How similar are two states  $\rho$  and  $\sigma$ .
- Defined by  $F(\rho, \sigma) := \text{Tr}(\sqrt{\sqrt{\rho}\sigma\sqrt{\rho}})^2$ .

## Average gate fidelity of a channel $\mathcal{E}$

$$\begin{aligned}\bar{F}(\mathcal{E}) &:= \int_{|\psi\rangle \in \mu} d\mu F(U|\psi\rangle\langle\psi|U^\dagger, \mathcal{E}(|\psi\rangle\langle\psi|)), \\ &= \int_{|\psi\rangle \in \mu} d\mu F(|\psi\rangle\langle\psi|, U^\dagger\mathcal{E}(|\psi\rangle\langle\psi|)U).\end{aligned}$$



# Introduction to Randomised Benchmarking

## Summary of randomised benchmarking

- Goal of randomised benchmarking: estimate  $\bar{F}$ .
- RB requires: a gate set  $\mathcal{G}$ , preparation of the state  $|0\rangle$ , a measurement  $\langle 0|$ .

## Group properties

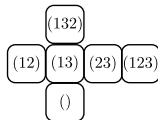
- Composition of two gates is another gate from the gate set.
- Each gate has an inverse operation.

Group elements  
 $g \in G$

Physical implementation or gate

$$\boxed{g} = \boxed{U_g}$$

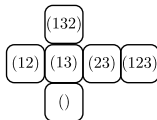
# Algorithm: toy example $S_3$



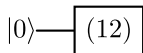
Prepare the state  $|0\rangle$

$|0\rangle$

# Algorithm: toy example $S_3$

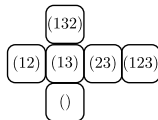


Roll the dice and apply the corresponding gate





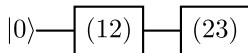
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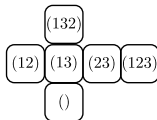
Roll a second time and apply the corresponding gate

(12)

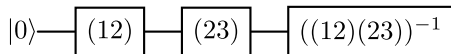
(23)



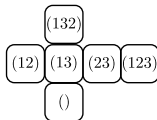
# Algorithm: toy example $S_3$



Apply the **inversion gate**



# Algorithm: toy example $S_3$

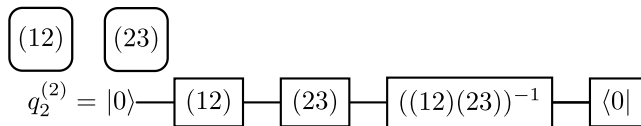
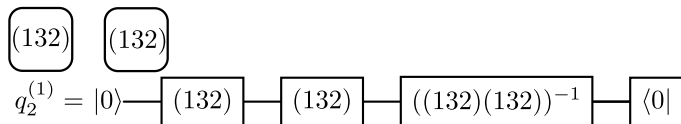


Measure with respect to the ground state  $\langle 0|$

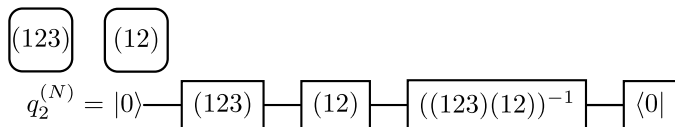
$$q_2^{(2)} = |0\rangle \text{---} \boxed{(12)} \text{---} \boxed{(23)} \text{---} \boxed{((12)(23))^{-1}} \text{---} \boxed{\langle 0|}$$

# Algorithm: toy example $S_3$

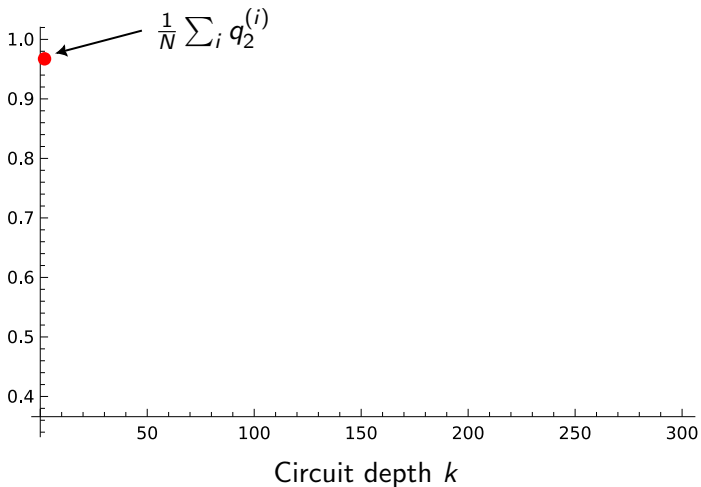
We repeat the process above  $N$  times: sample  $N$  sequences of 2 gates



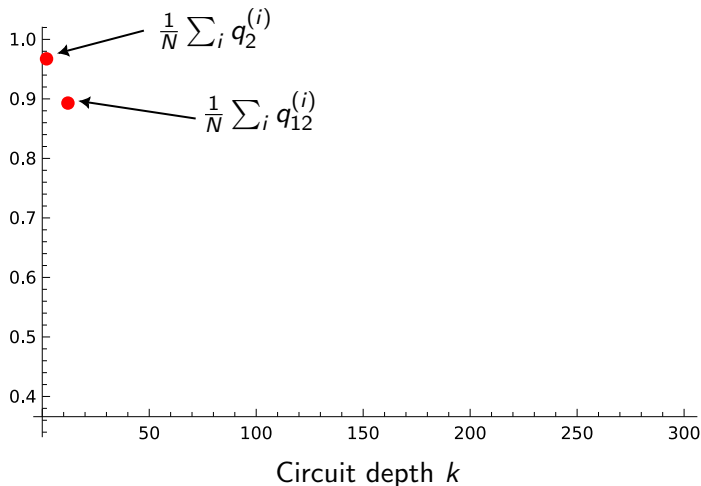
$\vdots$



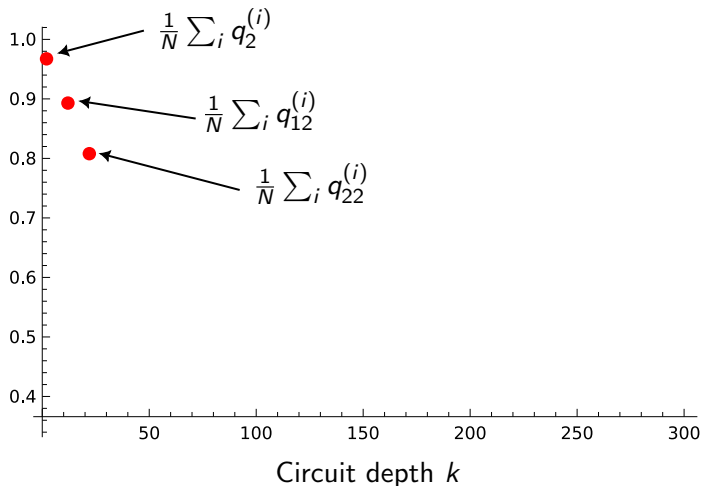
# Survival probability curve; Clifford



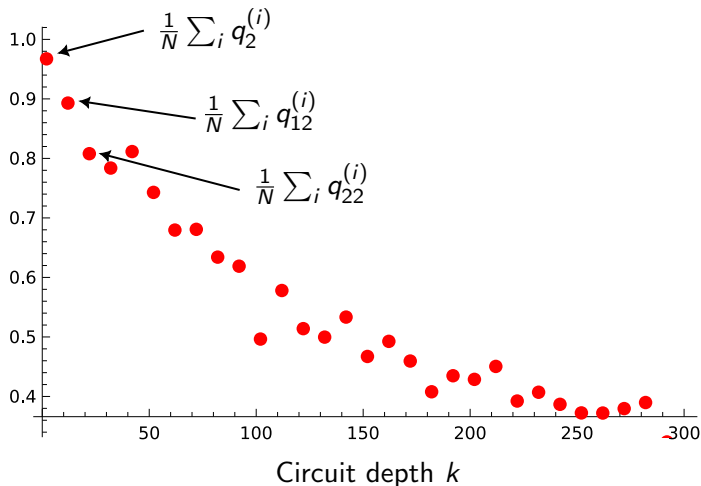
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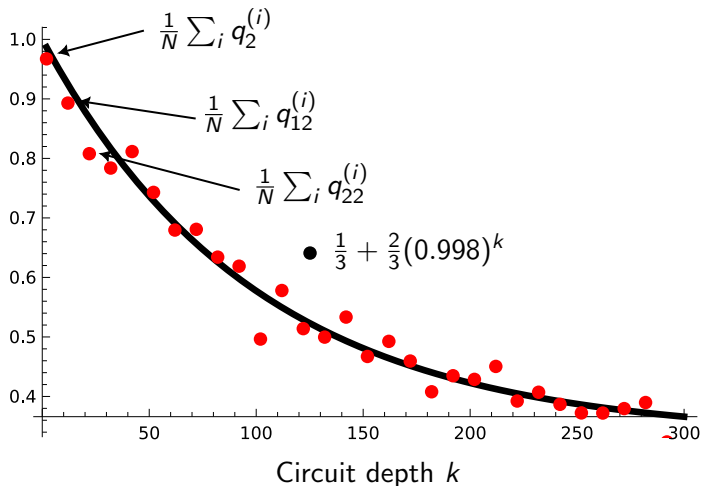


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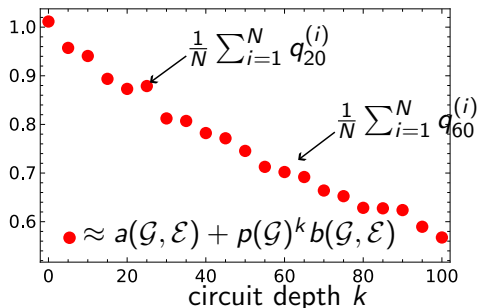




# Survival probability curve; Clifford



- Survival probability curve.

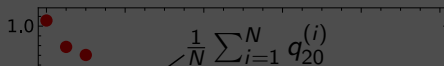


- It can be proven with some representation theory results that:

$$\bar{F}(\mathcal{G}) = \frac{3p(\mathcal{G}) - 1}{2}. \quad (1)$$

- (Recall that)  $\bar{F}$  is the average fidelity over a gate set.

- Survival probability curve.



## Attention

- Randomised benchmarking estimates  $\bar{F}$  by fitting a single exponential curve to the survival probability.

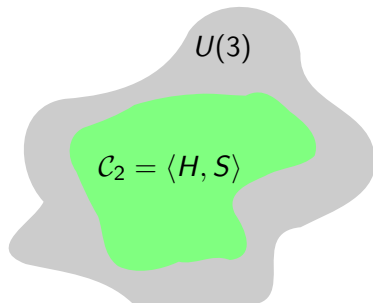
$$\bar{F}(\mathcal{G}) = \frac{3\rho(\mathcal{G}) - 1}{2}. \quad (1)$$

- (Recall that)  $\bar{F}$  is the average fidelity over a gate set.

# Clifford is not universal

## Clifford and universality

- Can be efficiently simulated on a classical computer.
- **Cannot** efficiently approximate an arbitrary unitary.



$$H := 3^{-1/2} \begin{bmatrix} 1 & 1 & 1 \\ 1 & \omega & \omega^2 \\ 1 & \omega^2 & \omega \end{bmatrix}, S := \begin{bmatrix} 1 & & \\ & 1 & \\ & & \omega^2 \end{bmatrix}. \quad (2)$$

# Universality $\Rightarrow$ T gates

## How do we achieve universality?

For qutrits, only two gates are necessary to approximate an arbitrary unitary operation: H (which is Clifford) and T.

## T gate

- A T gate is any gate that is a generator of a universal gate set which is non-Clifford.
- Conjugating Pauli matrices by a T gate produces a Clifford gate.
- For instance,

$$T = \begin{bmatrix} 1 & & \\ & \exp(2\pi i/9) & \\ & & \exp(2\pi i/9)^8 \end{bmatrix}; \quad T^9 = \mathbb{I}. \quad (3)$$

# The HDG: characterisation of $T$ (my contribution)

HDG stands for Hyperdihedral group.

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## Motivation

- Generalise dihedral benchmarking [1].
- Fewer group elements than the Clifford group.
- Smallest number of parameters in the expression for  $\bar{F}$ .

# The HDG: characterisation of $T$ (my contribution)

HDG stands for Hyperdihedral group.

## Motivation

- Generalise dihedral benchmarking [1].
- Fewer group elements than the Clifford group.
- Smallest number of parameters in the expression for  $\bar{F}$ .

## Construction

- Every product of  $X$  and  $T$ .

$$X := \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}, T = \begin{bmatrix} 1 & & \\ & \exp(2\pi i/3) & \\ & & \exp(2\pi i/3)^8 \end{bmatrix}. \quad (4)$$

- $X^x T^\alpha T'^\beta$ ,  $x \in [2] := \{0, 1, 2\}$  and  $\alpha, \beta \in [8]$ .
- $T' := XTX^\dagger$ .

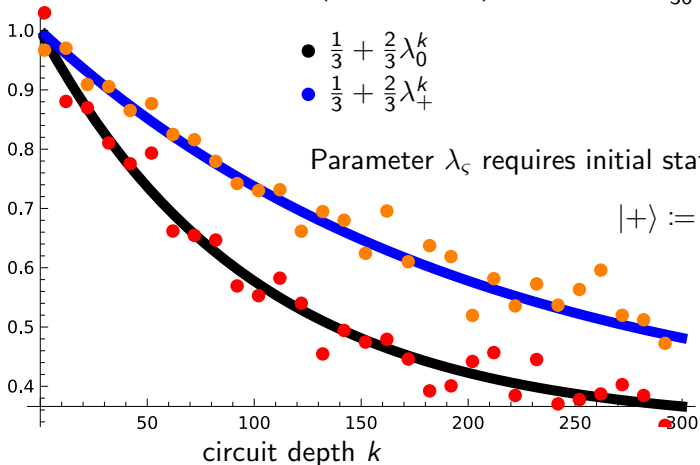


# Resulting scheme: HDG randomised benchmarking

The average gate fidelity (for the HDG) is  $\bar{F} = \frac{9+3(1+2\lambda_0+6\lambda_+)}{36}$ .

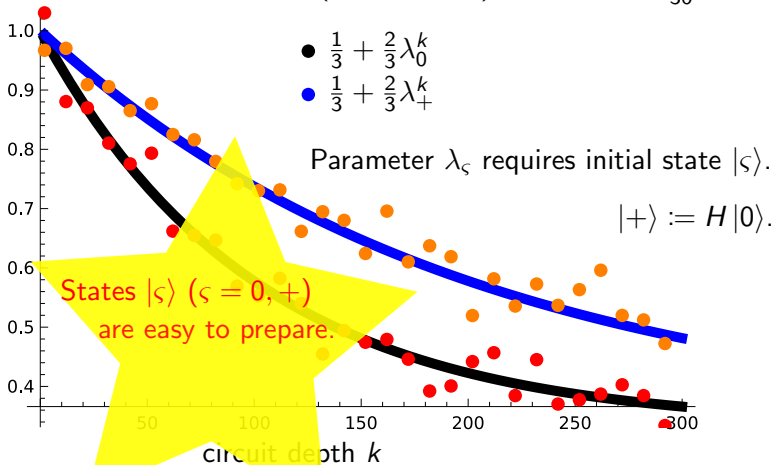
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# Conclusions

- Estimating the performance of quantum gates using the average gate fidelity.
- Randomised benchmarking is the Academy and industry standard for the performance quantification of quantum gates.
- I presented my generalisation, mainly done by the introduction of a group HDG, of the randomised benchmarking scheme to characterise universal qutrits gates.



## Collaborators

- Dr. Hubert de Guise (Lakehead University).
- Dr. Barry C. Sanders (University of Calgary).

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# Why two parameters?

- $\Gamma(\text{HDG}) = \Sigma_I + \Sigma_0 + \Sigma_0^* + \Sigma_+ + \Sigma_+^*$ .
- Each irrep has associated a parameters in the average gate fidelity.
- For high-fidelity gates, the phase is negligible.

## PL representation [8]

- Given an orthonormal basis for the set of  $3 \times 3$  matrices:  $\{W_i\}$ .
- $\Gamma(\mathcal{E})_{i,j} := 3^{-1/2} \text{Tr}(W_i^\dagger \mathcal{E}(W_j))$ .
- $\Gamma$  is a representation.

## Irreps

For a unitary irrep  $\gamma$ ,  $\Gamma \cong \gamma \otimes \bar{\gamma}$ , where  $\bar{\cdot}$  denotes complex conjugate.

## Average gate fidelity [11]

For a given channel  $\mathcal{E}$ ,

$$\bar{F}(\mathcal{E}, \mathbb{I}) = \frac{3 \text{Tr}(\Gamma) + 9}{36}. \quad (5)$$

If the noise  $\mathcal{E}$  is the same for each gate-set member  $\bar{F}(\mathcal{E}, \mathbb{I}) = \bar{F}(\mathcal{G})$ .

# Survival probability theoretical curve

It can be shown that:

$$\sum_{\mathbf{g} \in G^k} \langle 0 | \tilde{U}_{\text{inv}(\mathbf{g})} \tilde{U}_{g_1} \cdots \tilde{U}_{g_k} | 0 \rangle = \langle\langle 0 | \mathcal{T}(\mathcal{E})^k | 0 \rangle\rangle, \quad (6)$$

where  $\tilde{U} := \rho \mapsto \mathcal{E}(U\rho U^\dagger)$  and

$$\mathcal{T}(\mathcal{E}) := \sum_{g \in G} \Gamma(U_g)^\dagger \Gamma(\mathcal{E}) \Gamma(U_g). \quad (7)$$

The matrix  $\mathcal{T}$  is known as the **twirl** of  $\mathcal{E}$ .

## Normaliser

Let  $N$  and  $G$  groups such that  $N \leq G$ .  $G$  **normalises**  $N$  if  $\forall g \in G, h \in N, ghg^{-1} \in N$ .

## Pauli group for qutrits

$X$  and  $Z$  matrices and phases  $\omega := \exp(2\pi i/3)$ .

## Characteristics

- **Normaliser of the Pauli** group.
- Can be efficiently simulated on a classical computer.
- **Cannot** efficiently **approximate an arbitrary unitary**.

# Universal gates

## Universal gates

- Gates with up to three qutrits.
- Finite set of gates approximate up to arbitrary accuracy any evolution.

## Efficient

- Polylog with respect to the inverse of the accuracy  $\varepsilon$ .
- Meaning, increasing one digit of precision in accuracy increases polynomially.

## Solovay-Kitaev theorem [12]

- Efficient computation of the circuit.
- Efficient circuit depth.