

Neural Information Causality

A Random-Access Capacity Law for Representation Learning

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Core thesis

A query-separated representation is operationally a message.

- ▶ The encoder commits to a data-dependent interface before the query is known.
- ▶ The decoder later answers using only that interface, the query, and allowed side resources.
- ▶ Therefore random-access information must be bounded by independently certified capacity.

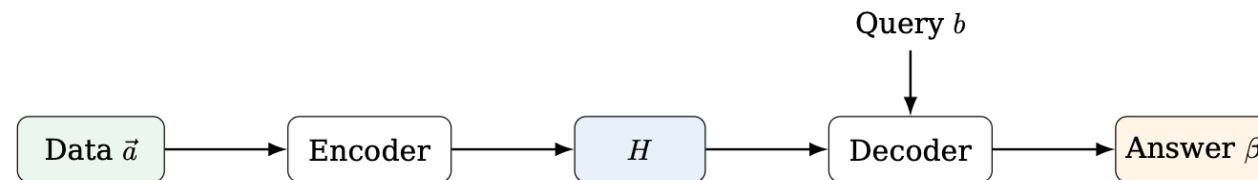
$$I_{\text{N-RAC}} \leq I(\vec{a} : H, B) \leq C_H$$

Why this question matters

Modern learning systems often rely on internal interfaces rather than direct data access.

- ▶ Memory-augmented networks and differentiable memory.
- ▶ Retrieval-augmented generation and query-time retrieval.
- ▶ Encoder–decoder pipelines with a hidden representation.
- ▶ Quantum- or correlation-assisted modules used inside learning architectures.

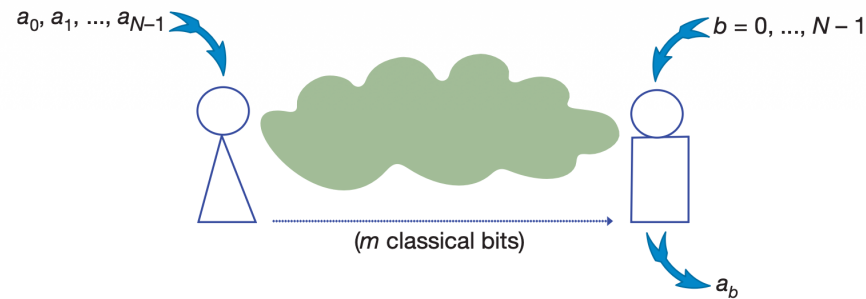
The diagnostic question: does the representation carry only its counted capacity?



***Query-separated structure**

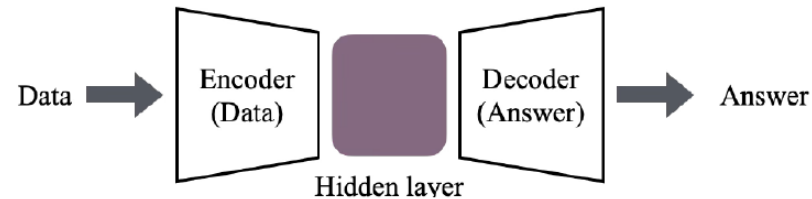
The shared causal skeleton

(a) Information Causality



M. Pawłowski et al., Nature 461, 1101 (2009)

(b) Deep Learning: Representation Learning



J.Bang and M. Pawłowski, arXiv:2605.09316 (2026)

- ▶ IC/RAC: Alice compresses data into an m -bit message.
- ▶ Representation learning: the hidden layer plays the same causal role.

Classical information causality: RAC form

- ▶ Alice receives N independent unbiased bits:

$$\vec{a} = (a_0, \dots, a_{N-1}).$$

- ▶ Bob receives a uniformly random index b .
- ▶ Alice sends $x \in \{0, 1\}^m$.
- ▶ Bob outputs β as a guess for a_b .

$$P_K = \Pr[\beta = a_K \mid b = K]$$

$$I_{\text{RAC}} = \sum_{K=0}^{N-1} I(a_K : \beta \mid b = K)$$

Information causality: $I_{\text{RAC}} \leq m$.

Accuracy is not the score, but it certifies score

The observable information score uses only classical variables $(\vec{a}, b, \beta)$:

$$I_{\text{RAC}} = \sum_K I(a_K : \beta \mid b = K).$$

For unbiased target bits, binary estimation gives

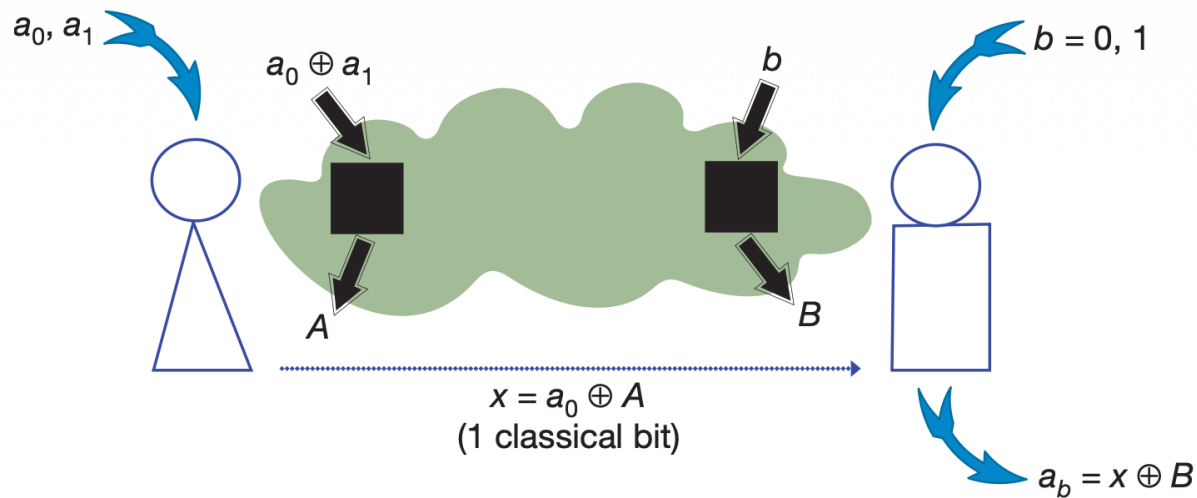
$$I(a_K : \beta \mid b = K) \geq 1 - h(P_K).$$

Therefore

$$I_{\text{RAC}} \geq N - \sum_{K=0}^{N-1} h(P_K), \quad I_{\text{RAC}} \geq N[1 - h(P)] \text{ if } P_K \equiv P.$$

This converts random-access accuracy into an information-accounting test.

The seed case: $N = 2$, $m = 1$, and CHSH

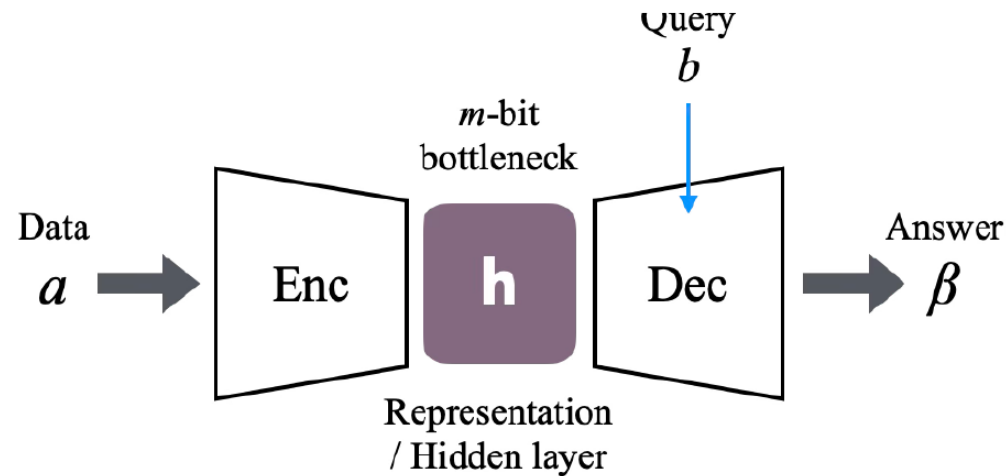


- ▶ Alice uses $s = a_0 \oplus a_1$.
- ▶ Bob uses $t = b$.
- ▶ The correlation cell outputs (A, B) .
- ▶ Message and answer:

$$x = a_0 \oplus A, \quad \beta = x \oplus B.$$

Correct retrieval is exactly the CHSH winning event.

Neural-RAC: the query-separated primitive



- ▶ The encoder is query-blind.
- ▶ The decoder sees the query only after H is fixed.
- ▶ W is fixed before the fresh episode.
- ▶ H is the only data-dependent causal carrier.

$$H \leftarrow \text{Enc}(\vec{a}; R_A, W), \quad \beta \leftarrow \text{Dec}(H, b; R_B, W).$$

Capacity must be certified independently

The effective bottleneck capacity is not inferred from success after the fact:

$$C_H(H) = \sup I(\vec{a} : H \mid B).$$

Useful certificates include:

- ▶ finite alphabet: $C_H \leq \log |\mathcal{H}|$;
- ▶ d coordinates quantized to q bits: $C_H \leq dq$;
- ▶ AWGN interface: $C_H \leq \frac{d}{2} \log(1 + P/\sigma^2)$;
- ▶ hard classical bottleneck: $H \in \{0, 1\}^m$, so $C_H \leq m$.

Unconstrained ideal real values do not imply a finite capacity law.

Main theorem: embedding and capacity

Every query-separated Neural-RAC architecture induces a random-access communication experiment.

$$I_{\text{N-RAC}} \leq I(\vec{a} : H, B)$$

If an independent interface certificate gives $I(\vec{a} : H, B) \leq C_H$, then

$$I_{\text{N-RAC}} \leq C_H.$$

For a hard m -bit bottleneck:

$$I_{\text{N-RAC}} \leq m.$$

Neural-IC as a capacity audit

The theorem becomes an evaluation protocol:

1. Fix all parameters W before the episode.
2. Sample a fresh database $\vec{a}$ independently of W and the resource.
3. Encode H before observing the query b .
4. Decode β using only (H, b) and allowed side resources.
5. Estimate $I_{\text{N-RAC}}$ and compare with the certified C_H .

If $I_{\text{N-RAC}} > C_H$, the model has an uncounted side channel or an invalid capacity certificate.

Capacity–accuracy tradeoff

Combining the accuracy lower bound with Neural-IC gives

$$\text{Entropy defects } N - \sum_{K=0}^{N-1} h(P_K) \leq C_H. \rightarrow \text{a sanity check}$$

For query-symmetric access:

$$C_H \geq N[1 - h(P)].$$

A tiny representation with uniformly high random-access accuracy implies one of four diagnoses:

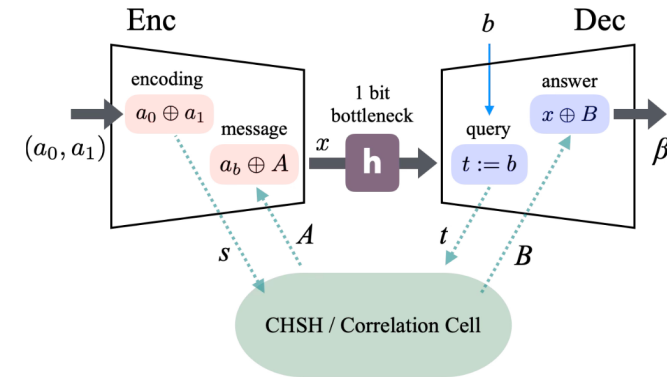
- ▶ query leakage into the encoder;
- ▶ finite-precision leakage through real-valued interfaces;
- ▶ episode-specific weights or memory;
- ▶ uncounted nonclassical or post-quantum resources.

A CHSH correlation cell as a neural resource layer

The shared resource is modeled as independent no-signaling CHSH cells:

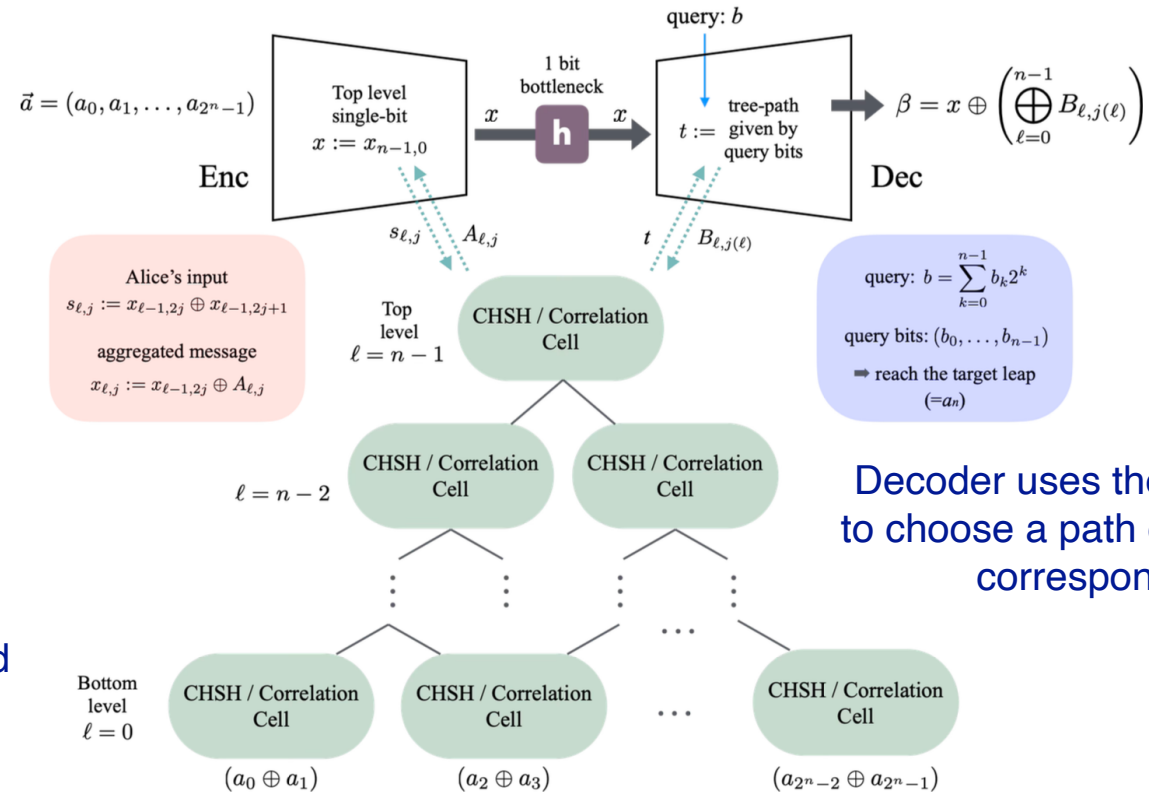
$$\Pr[A \oplus B = st \mid s, t] = \frac{1 + E}{2}.$$

- ▶ E is the correlation bias.
- ▶ Classical local resources satisfy $E \leq 1/2$.
- ▶ Quantum theory permits $E \leq 1/\sqrt{2}$.
- ▶ Post-quantum no-signaling boxes could have stronger E .



What happens when such a cell is used repeatedly inside a query-separated architecture.

Nesting: a depth- n pyramid for $N = 2^n$



Higher levels recursively combine those messages

pairs of database bits are combined into intermediate messages

Decoder uses the binary expansion of the query to choose a path down the tree and combines the corresponding correlation outputs.

Alice aggregates to one bit x ; Bob follows the query path through independent correlation cells.

Bias multiplication is the amplification mechanism

Each selected CHSH cell contributes an independent error bit. If local correctness bias is E , the parity of n errors has bias E^n .

$$P_K = \Pr[\beta = a_K \mid b = K] = \frac{1 + E^n}{2} \quad \text{*bias multiplication}$$

*Along a fixed query path, each CHSH cell contributes an independent error bit.

For $N = 2^n$:

$$I_{\text{N-RAC}}(n, E) = 2^n \left[1 - h\left(\frac{1 + E^n}{2}\right) \right].$$

Small single-cell advantages can become a depth-dependent random-access amplifier.

Tsirelson threshold from one-bit stability

The entropy deficit obeys

$$1 - h\left(\frac{1 + \delta}{2}\right) \geq \frac{\delta^2}{2 \ln 2}.$$

*quadratic lower bound

With $\delta = E^n$ and $N = 2^n$:

$$I_{\text{N-RAC}} \geq \frac{(2E^2)^n}{2 \ln 2}.$$

- ▶ If $E > 1/\sqrt{2}$, then $2E^2 > 1$ and the score eventually exceeds one bit.
- ▶ Stability of a one-bit bottleneck for arbitrary depth therefore requires $E \leq 1/\sqrt{2}$.

Tsirelson's value appears as the stability frontier for representation-mediated random access.

Beyond isotropy: asymmetric branch biases

If the seed cell has different branch biases E_0 and E_1 , a query path $b = (b_1, \dots, b_n)$ gives

$$P_b = \frac{1 + \prod_{\ell=1}^n E_{b_\ell}}{2}.$$

The exact score for binary-symmetric branches is

$$\sum_{b \in \{0,1\}^n} \left[1 - h \left(\frac{1 + \prod_{\ell} E_{b_\ell}}{2} \right) \right].$$

A sufficient amplification criterion is

$$E_0^2 + E_1^2 > 1.$$

This keeps task-specific asymmetry visible instead of collapsing everything into one scalar too early.

Why quantum resources obey the accounting law

Quantum mutual information satisfies the information-calculus assumptions needed in the theorem:

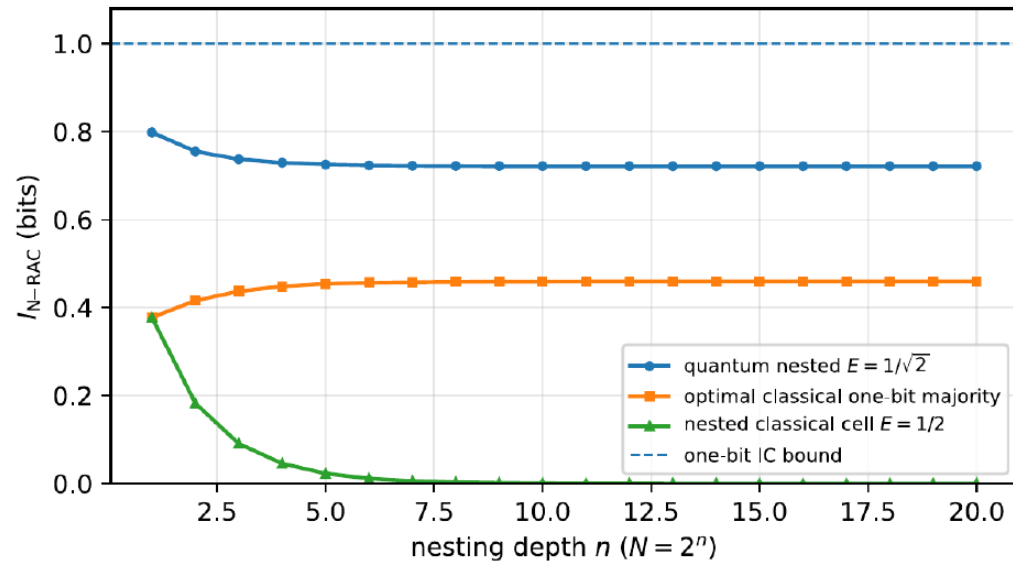
- ▶ consistency with Shannon mutual information on classical registers;
- ▶ product-state normalization;
- ▶ data processing under local CPTP maps;
- ▶ chain rule and nonnegative conditional mutual information;
- ▶ classical self-conditioning.

Therefore, quantum-assisted Neural-RAC protocols obey

$$I_{\text{N-RAC}} \leq m$$

for a hard classical m -bit bottleneck, and $I_{\text{N-RAC}} \leq C_H$ for certified interfaces.

Classical one-bit fair-access benchmark

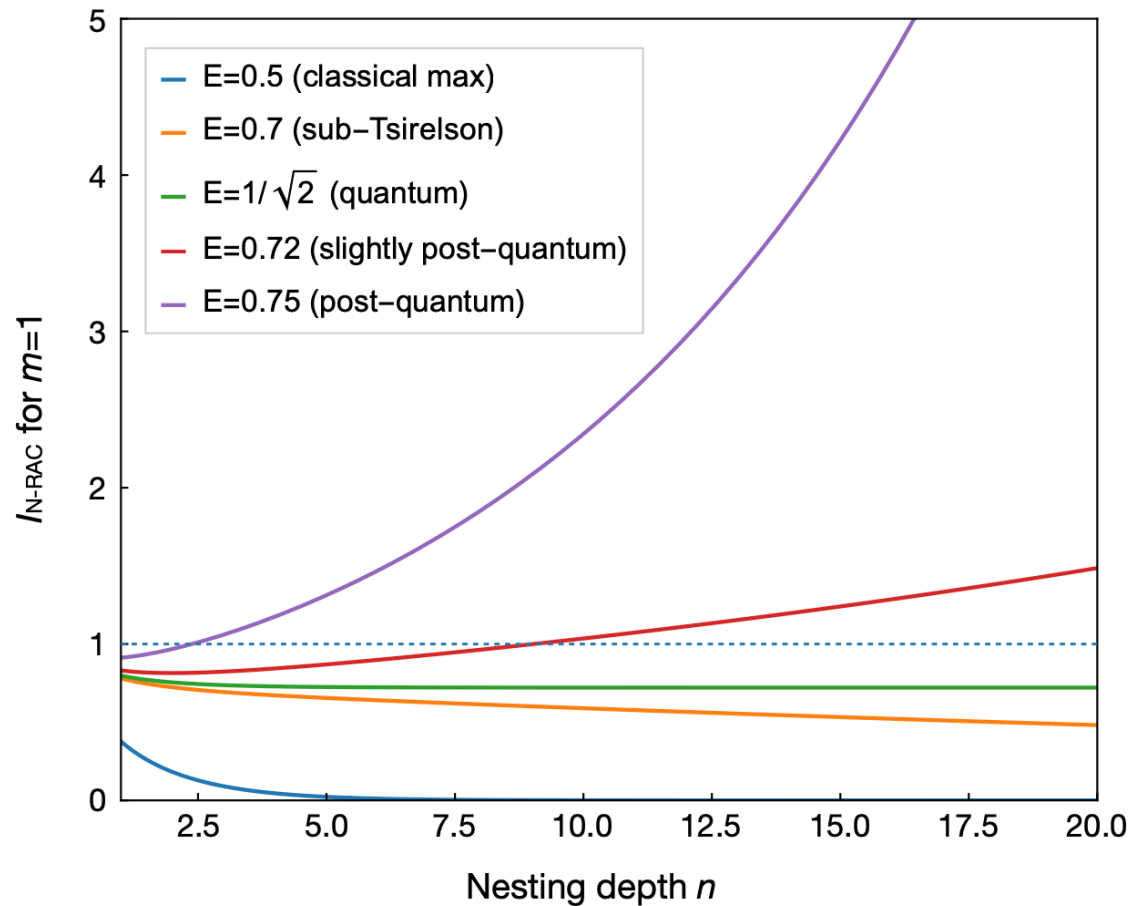


- ▶ One classical bit can saturate total information by storing one target bit.
- ▶ Fair access asks for average performance over all queries.

$$P_{\text{avg}}^{\text{cl}}(N, 1) = \frac{1}{2} + 2^{-N} \binom{N-1}{\lfloor (N-1)/2 \rfloor}$$

Quantum correlations improve fair access without exceeding the one-bit information budget.

Closed-form stability probe: score versus depth

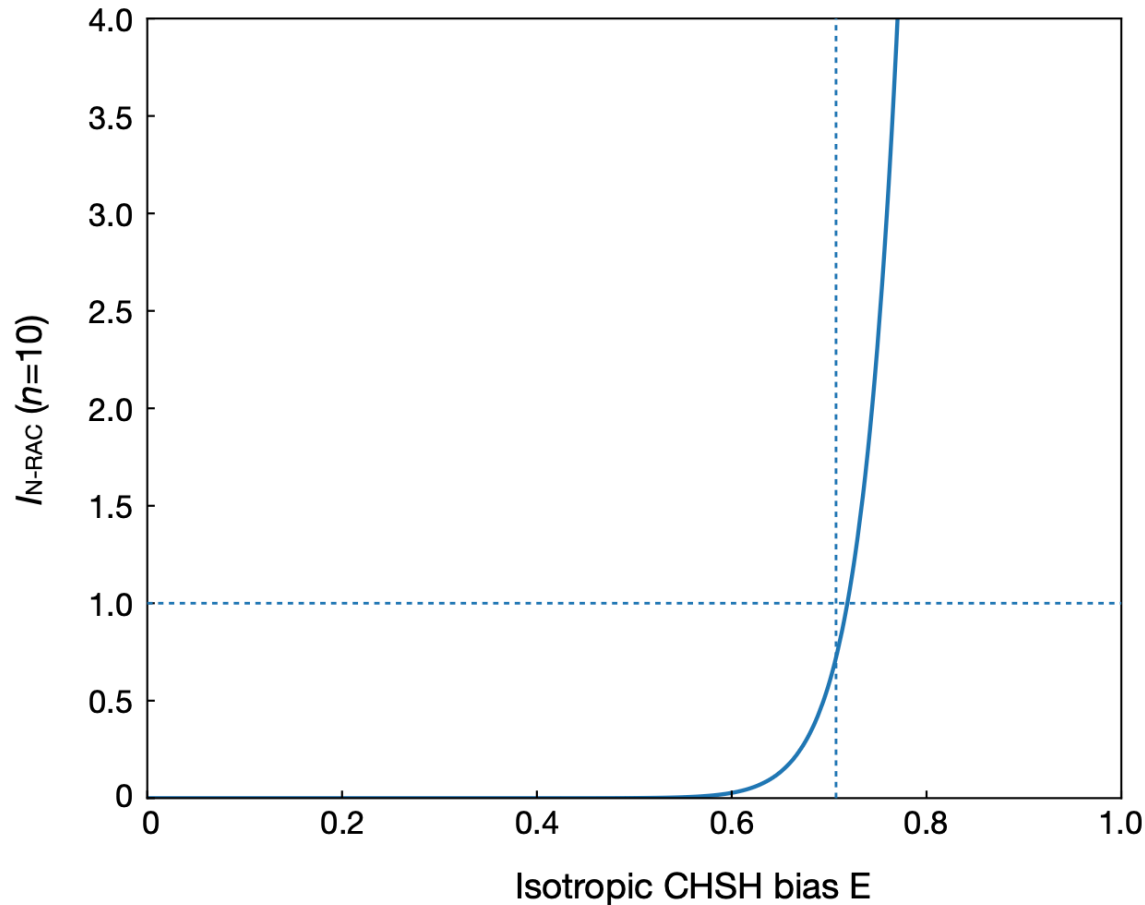


- ▶ Sub-Tsirelson curves remain stable.
- ▶ The quantum curve approaches a constant below one.
- ▶ Supercritical curves eventually cross the one-bit bound.

$$\lim_{n \rightarrow \infty} I_{N-RAC} \left(n, \frac{1}{\sqrt{2}} \right) = \frac{1}{2 \ln 2}$$

$$\frac{1}{2 \ln 2} \simeq 0.721.$$

Bias scan: shallow tests can hide supercriticality



At fixed depth $n = 10$, the finite-depth crossing is above the asymptotic threshold.

n	0.5	0.7	$1/\sqrt{2}$	0.72	0.75
1	.377	.780	.798	.832	.913
5	.023	.655	.725	.870	1.312
10	.001	.589	.721	1.036	2.344
20	.000	.482	.721	1.486	7.607

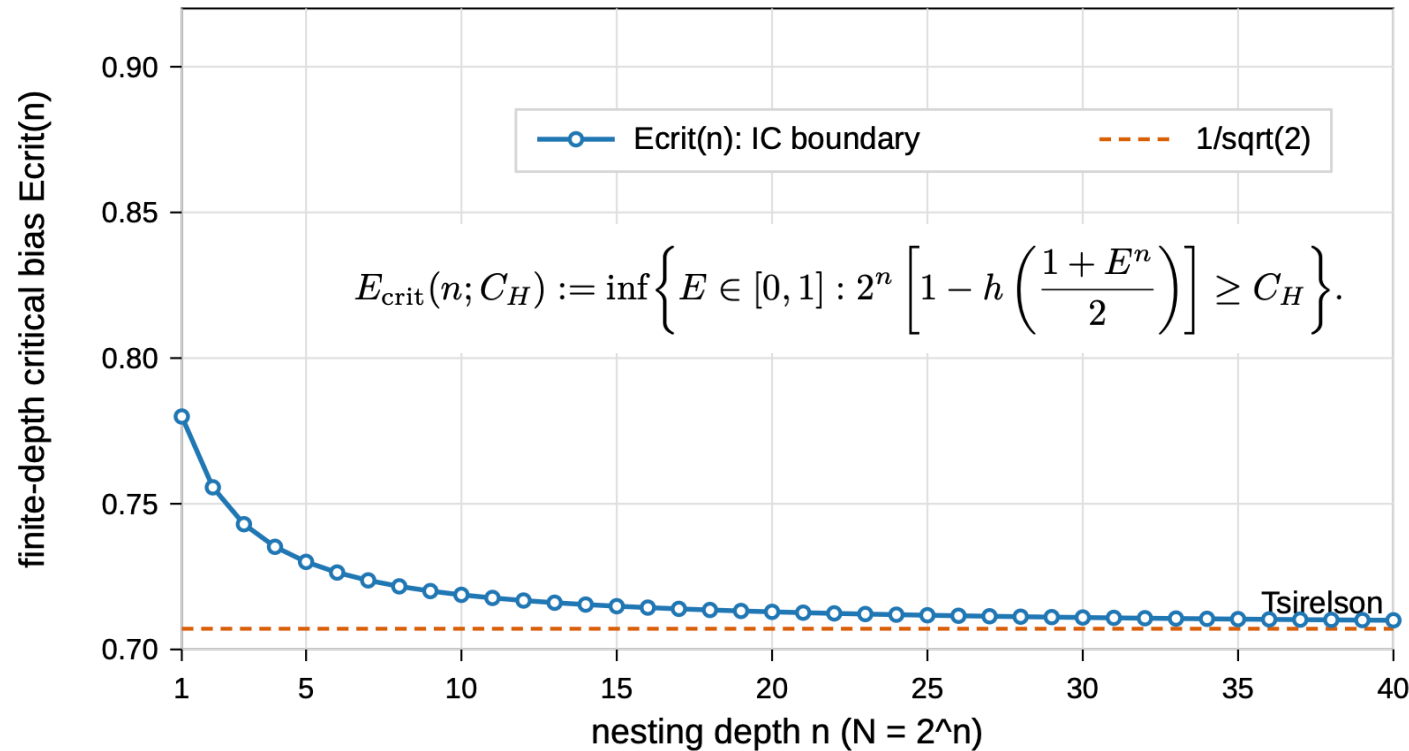
Values are $I_{N-RAC}(n, E)$; the one-bit limit is 1.

The vertical dashed line: Tsirelson value

The horizontal dashed line: Neural IC limit

The finite-depth tests have their own threshold.

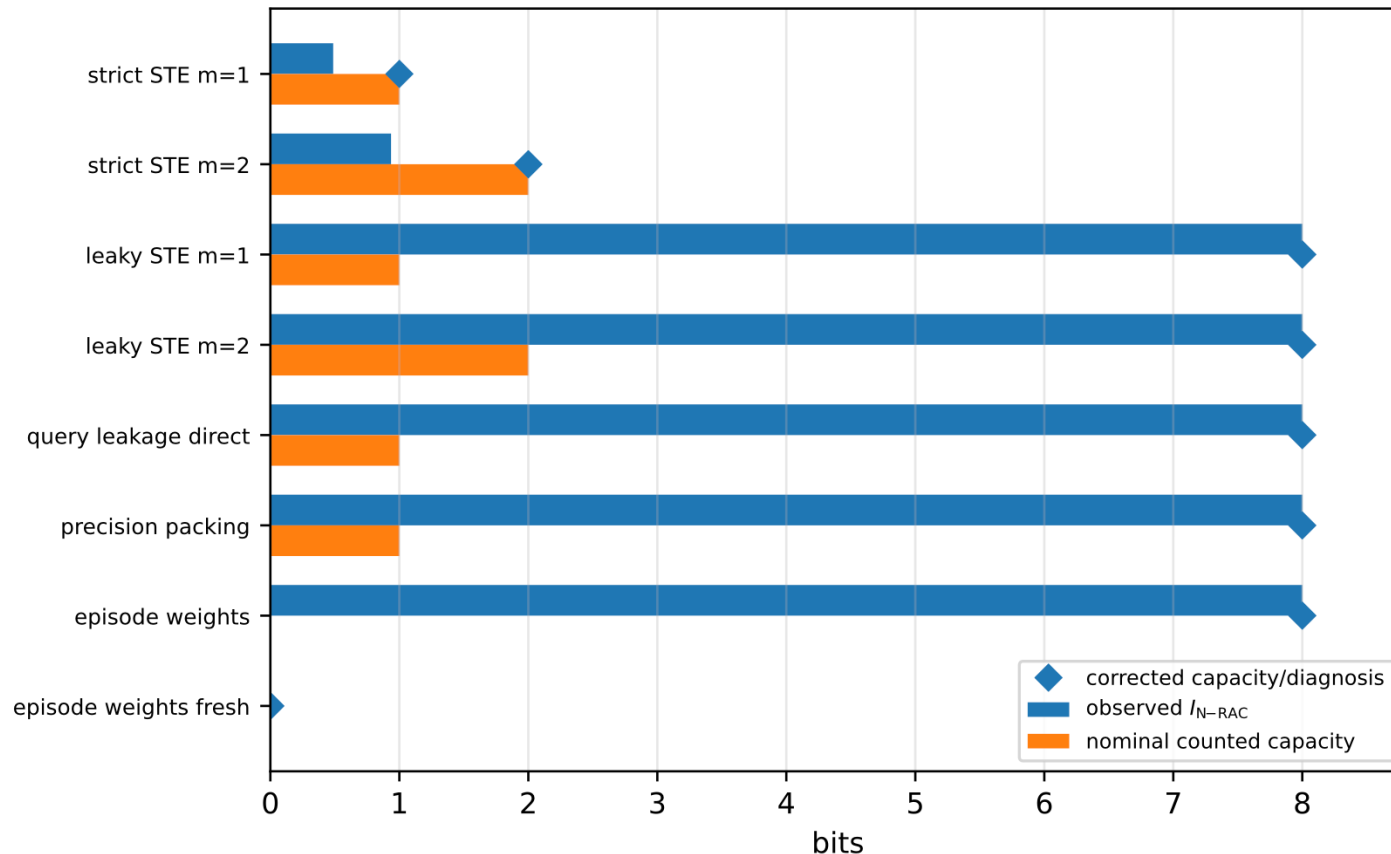
Finite-depth critical boundary



- $E_{\text{crit}}(n)$ solves $I_{\text{N-RAC}}(n, E) = 1$.
- It approaches $1/\sqrt{2}$ from above.
- Slightly post-quantum correlations may need sufficient depth before violation is visible.

The operational boundary tightens and approaches the Tsirelson value from above

Leakage diagnostics: strict versus leaky bottlenecks



- ▶ Strict binary bottlenecks stay below counted capacity.
- ▶ Query leakage can directly transmit a_b .
- ▶ Precision packing hides many bits in nominally small real vectors.
- ▶ Episode-specific weights behave as uncounted memory.

Neural-IC therefore gives a structured way to debug unusual powerful representations.

Take-home messages

1. Neural-IC is an operational embedding of IC into query-separated representation learning.
2. The key law is not dimension-based; it is capacity-certified:

$$I_{\text{N-RAC}} \leq I(\vec{a} : H, B) \leq C_H.$$

3. CHSH nesting turns correlation bias into an amplification mechanism.
4. Arbitrary-depth one-bit stability selects the Tsirelson threshold $E = 1/\sqrt{2}$.
5. Quantum correlations improve fair query-conditioned access without turning finite representations into oracles.

Thank you
