

Optimized solver algorithms for computing $Tr(A^{-1})$

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Lattice 2026 // UMD, College Park, MD

Introduction

- Estimating the trace of the inverse of a large quark matrix is a longstanding problem in lattice QCD.
- The standard estimator is **Hutchinson's method**, which requires solving systems of linear equations.
- In this talk I will explore:
 - **New solver algorithms** for Hutchinson's method
 - **Testing results**
 - **Singular value deflation** (SVD) and **polynomial approximation of A^{-1}** for computing $Tr(A^{-1})$



I. Hutchinson's method & solving algorithms

Hutchinson's method

Hutchinson's method for computing the trace of A -inverse is given by

$$\text{Tr}(A^{-1}) \approx \frac{1}{N} \sum_i^N b_i^\dagger A^{-1} b_i$$

where b is a random noise vector (z-4 noise in this work). In practice, N is not fixed but is dependent upon the desired trace tolerance.

It is prohibitively expensive to directly compute A^{-1} directly, so we solve the system $Ax = b$. The trace then becomes

$$\text{Tr}(A^{-1}) \approx \frac{1}{N} \sum_i^N b_i^\dagger x_i$$

Which linear equations solver do we use?

Solving algorithms at a glance

BiCGStab

- 2 MVPs per iteration
- 2 dot products per iteration (extremely costly on a distributed system)
- Usually numerically stable and straightforward to implement
- Independent solve per noise vector

BiCG-QCD

- BiCG algorithm specialized for QCD matrices
- Exploits the symmetry $\gamma_5 M = M^\dagger \gamma_5$ to eliminate a matrix vector product—1 MVP per iteration
- Premature breakdown is possible (zero divisors)
- Independent solve per noise vector

Twin BiCG

- One base solve, apply parameters to all subsequent solves.
- 1 MPV per iteration after base solve
- Eliminates all dot products beyond the base solve
- Easily adaptable to batched solves/multiple RHS at a time

Twin BiCG

Main Idea

Twin method applied to Bi-Conjugate Gradient (BiCG) recycles the same parameters from an initial base solve for all subsequent solves.

Theory

- Approximate solution is in the Krylov subspace,

$$\hat{x} \in \mathcal{K}^m(b, A) = \text{span}\{b, Ab, A^2b, \dots, A^{m-1}b\} \implies \hat{x} = p(A)b$$

Then

$$r = b - A\hat{x} = \pi(A)b$$

- To solve $Ax^{(i)} = b^{(i)}$, the polynomial $\pi(A) = I - Ap(A)$ must be small along the whole spectrum of A . This is true for each system solved.
- Reuse the polynomial on future solves!

Benefits of Twin

- MVPs are effectively cut in half
- Dot products are eliminated beyond the base solve
 - Important for large scale problems on distributed memory where the relative cost of a DP is very high
- Each solve requires the same number of iterations
 - Easily adapted to a batched solve context (i.e. solve multiple RHS at once until convergence)—further improves efficiency
- Only one good solve is needed!
 - If the problem is very difficult, independent solves may yield several non-convergences. Twin circumvents this by picking a successful system and reusing the parameters.

Summary

$$\text{Tr}(A^{-1}) \approx \frac{1}{N} \sum_i^N b_i^\dagger A^{-1} b_i$$

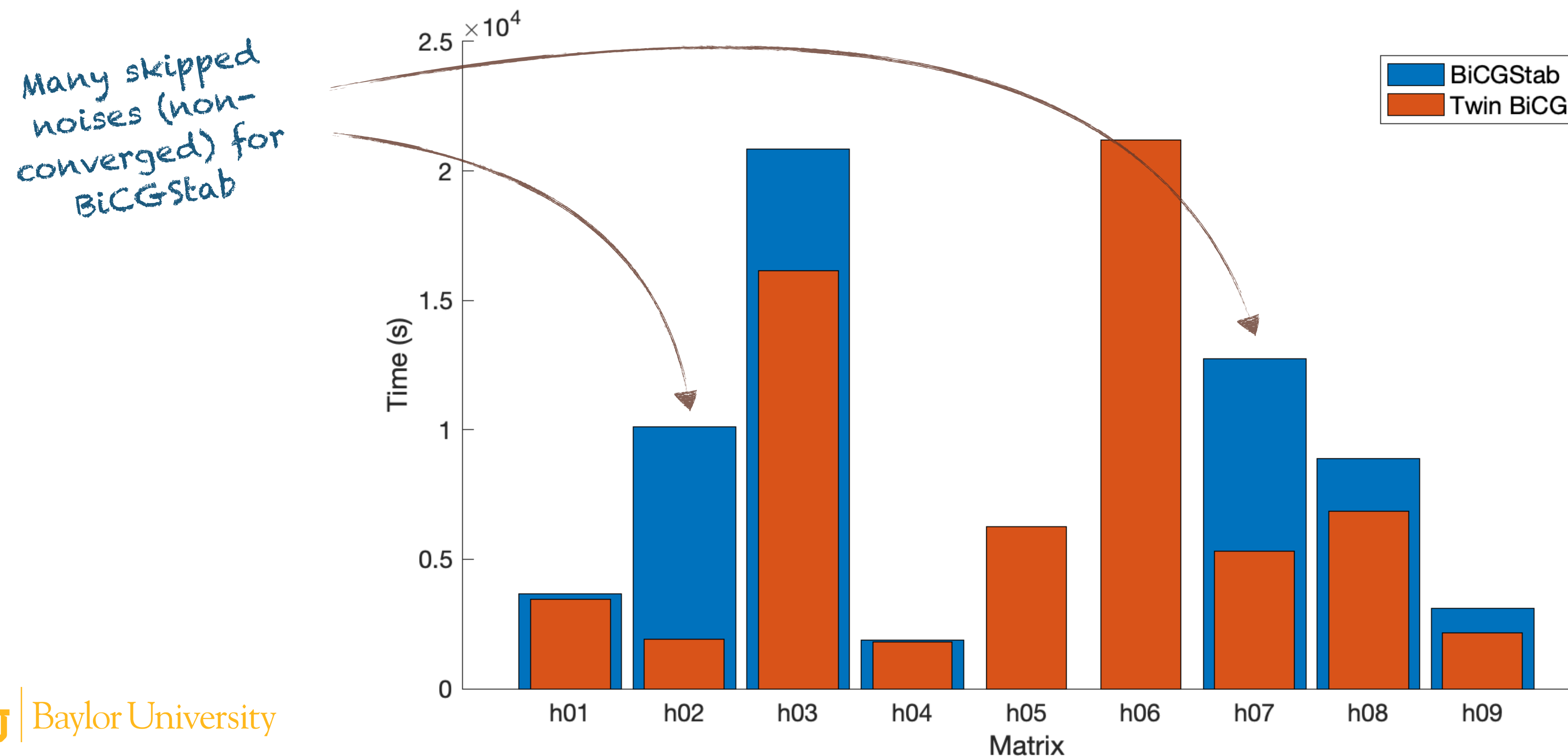
1. $n_{MVP} \xrightarrow{\text{Twin}} \frac{1}{2} n_{MVP}$
(after base solve)

2. $\sum_i n_{DP, i} = n_{DP, 1}$

3. $n_{iter, i>1} = n_{iter, i=1}$

Twin BiCG vs BiCGStab

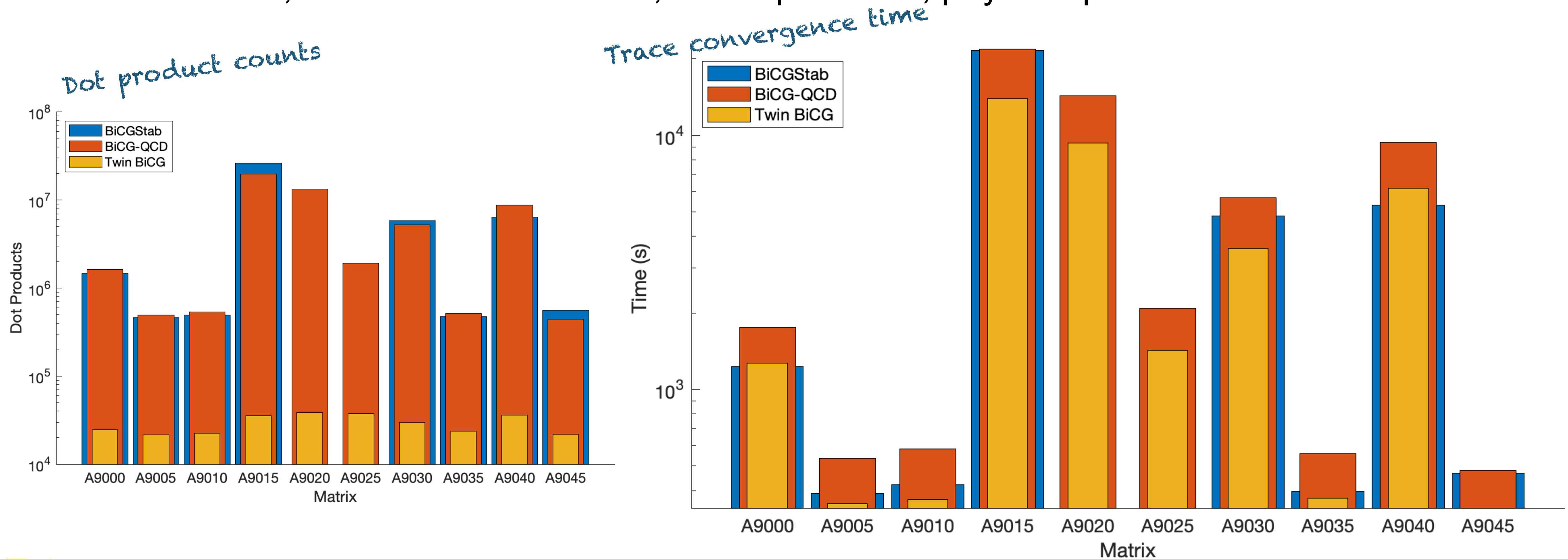
- Single thread, single CPU tests
- Directly compare algorithm performance of Twin BiCG against BiCGStab
- 9 matrices, $12^3 \times 16$ lattice, quenched, & κ_{crit}



Takeaway
Twin BiCG outperforms BiCGStab on the level of algorithm structure.

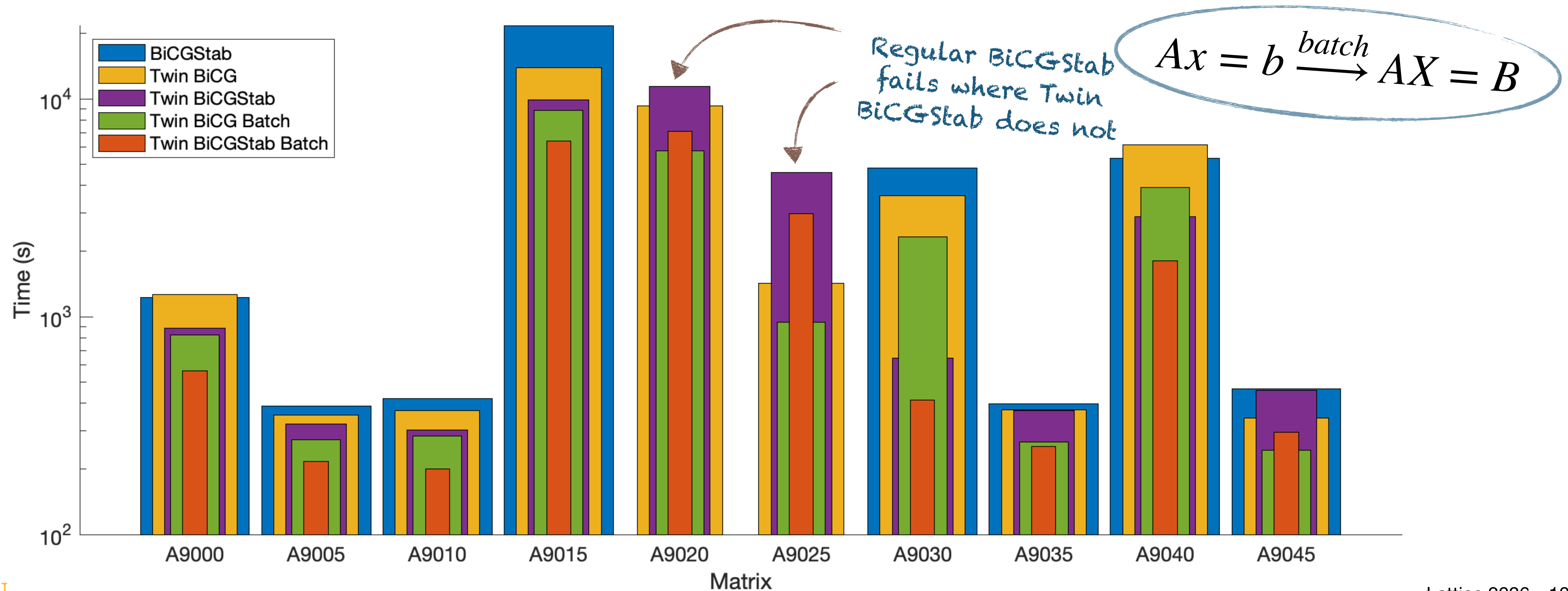
Twin BiCG, BiCGStab, BiCG-QCD

- Compare the three algorithms on NVIDIA A100 GPUs (and all tests hereafter)
- 10 matrices, $16^3 \times 48$ MILC lattice, semi-quenched, physical pion mass



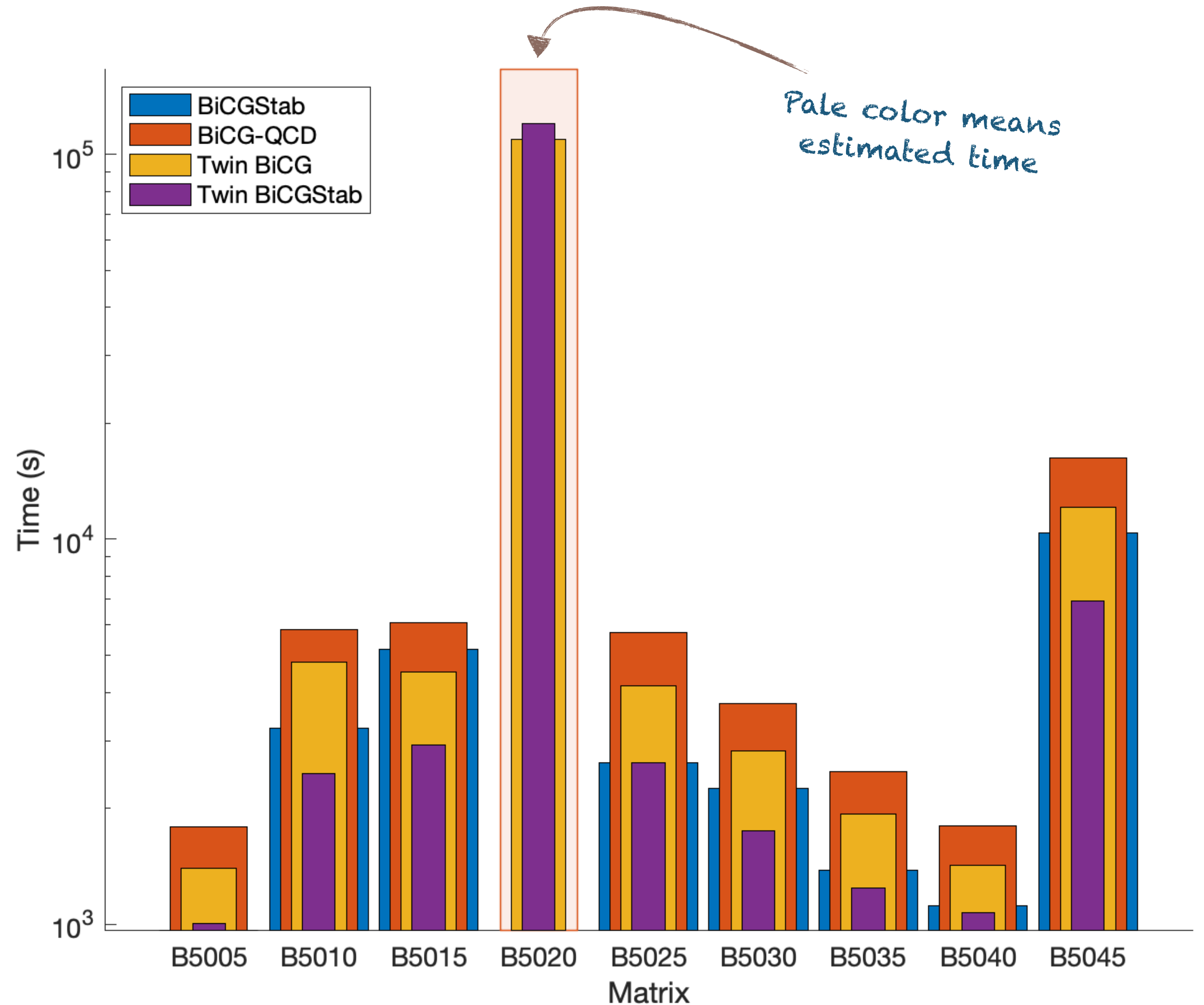
Include batching

- Same matrices: $16^3 \times 48$ MILC lattice size, semi-quenched, physical pion mass. Trace tol. $t_{tr} = 0.0005 \times 16^3 \times 48$, solve tol. $t_{solve} = 10^{-3}$
- **Batched** solves for Twin methods (BiCGStab included). (“Batching” solves multiple RHS at once.)

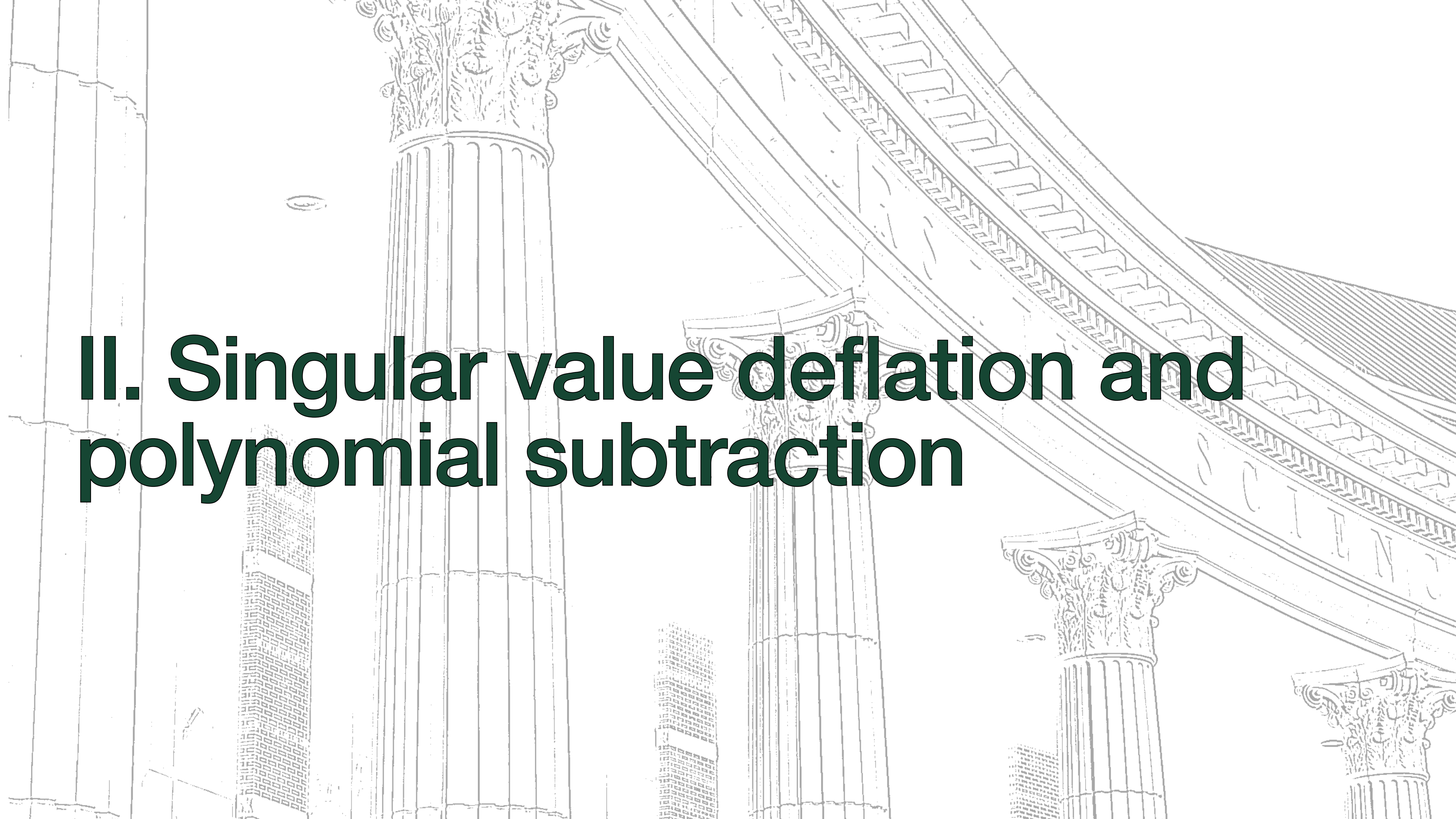


Larger matrix

- $24^3 \times 64$ MILC lattice, physical pion mass
- Trace tolerance
 $t_{tr} = 0.0005 \times 24^3 \times 64$
- Solve tolerance $t_{solve} = 10^{-3}$
- In instances of non-convergence, a good-faith time estimate is given, when possible.



II. Singular value deflation and polynomial subtraction



Polynomials and singular values

- The solver algorithms discussed stand in service of a more sophisticated evaluation of the trace,

$$Tr(A^{-1}) \approx \underbrace{Tr(A^{-1} - p(A))}_{(1)} + \underbrace{Tr\left(p(A) - \sum_i \frac{1}{\sigma_i} v_i u_i^\dagger\right)}_{(2)} + \underbrace{Tr\left(\sum_i \frac{1}{\sigma_i} v_i u_i^\dagger\right)}_{(3)}$$

- Polynomial $p(A) \approx A^{-1}$, vectors $\{v_i\}$ and $\{u_i\}$ are right and left singular vectors, respectively, $\{1/\sigma_i\}$ are the largest singular values of A^{-1} .
- When $p(A)$ is a strong approximation of A^{-1} , very few noises are needed for trace (1).
- Trace (2) deflates out the largest singular values, and trace (3) is computed exactly.
- $p(A)$ and SVD can be expensive to generate, but payoff is anticipated for difficult problems.

Polynomial approximation of A^{-1}

- Recall that the solution is extracted from the Krylov subspace

$$\hat{x} \in \mathcal{K}^m(b, A) = \text{span}\{b, Ab, A^2b, \dots, A^{m-1}b\} \implies \hat{x} = p(A)b$$

high-degree

- The polynomial $p(A)$ must approximate A^{-1} well
- Run non-symmetric Lanczos algorithm (specialized for QCD) on A
- In testing we focus on $\underbrace{Tr(A^{-1} - p(A))}_{(1)} + \underbrace{Tr(p(A))}_{(2)}$

Matrix	h01	h02	h03	h04	h05	h08	h09	h10
noises (1)	2	2	2	2	2	2	2	2
noises (2)	323	114	585	235	74	145	128	59

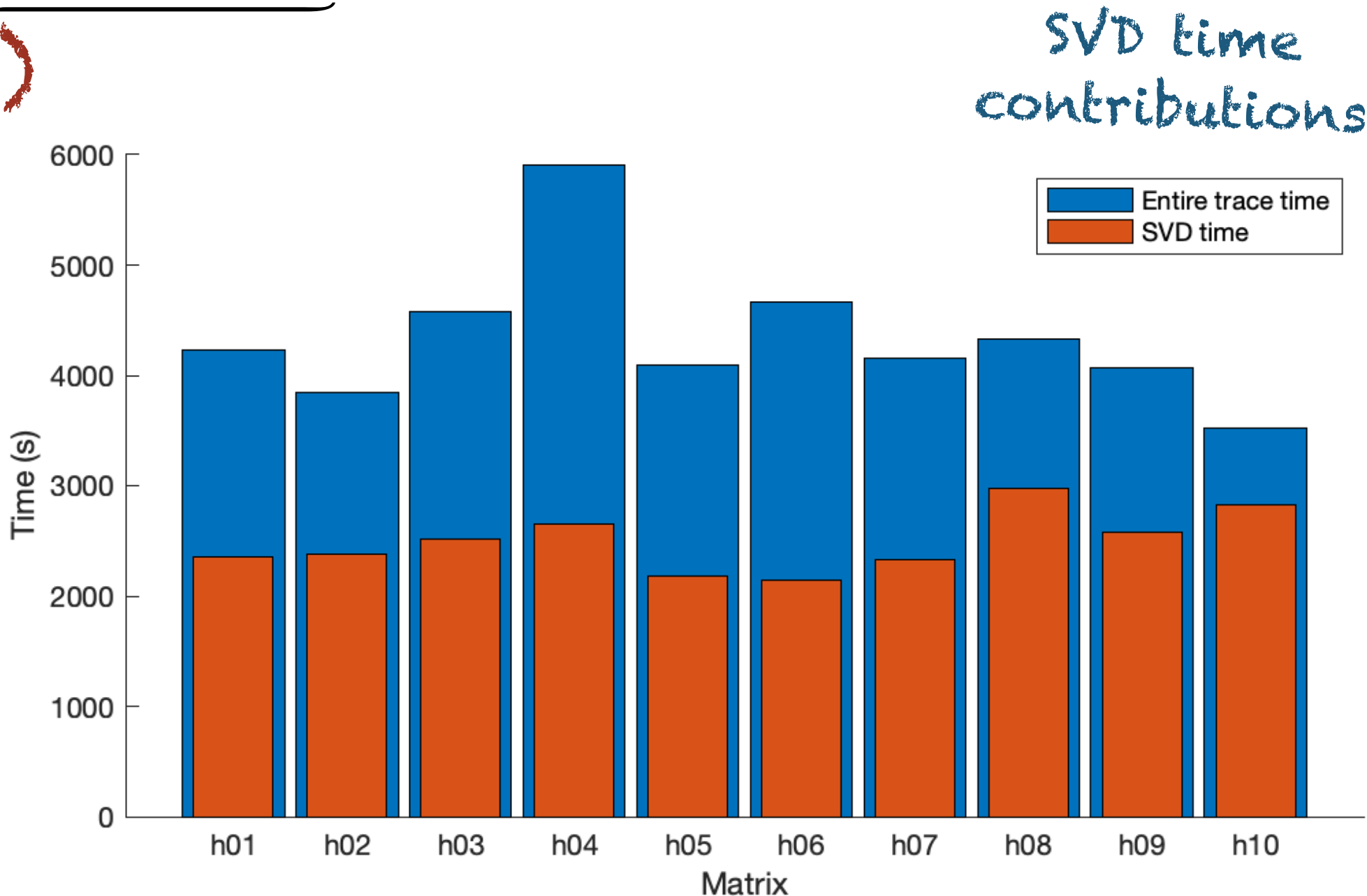
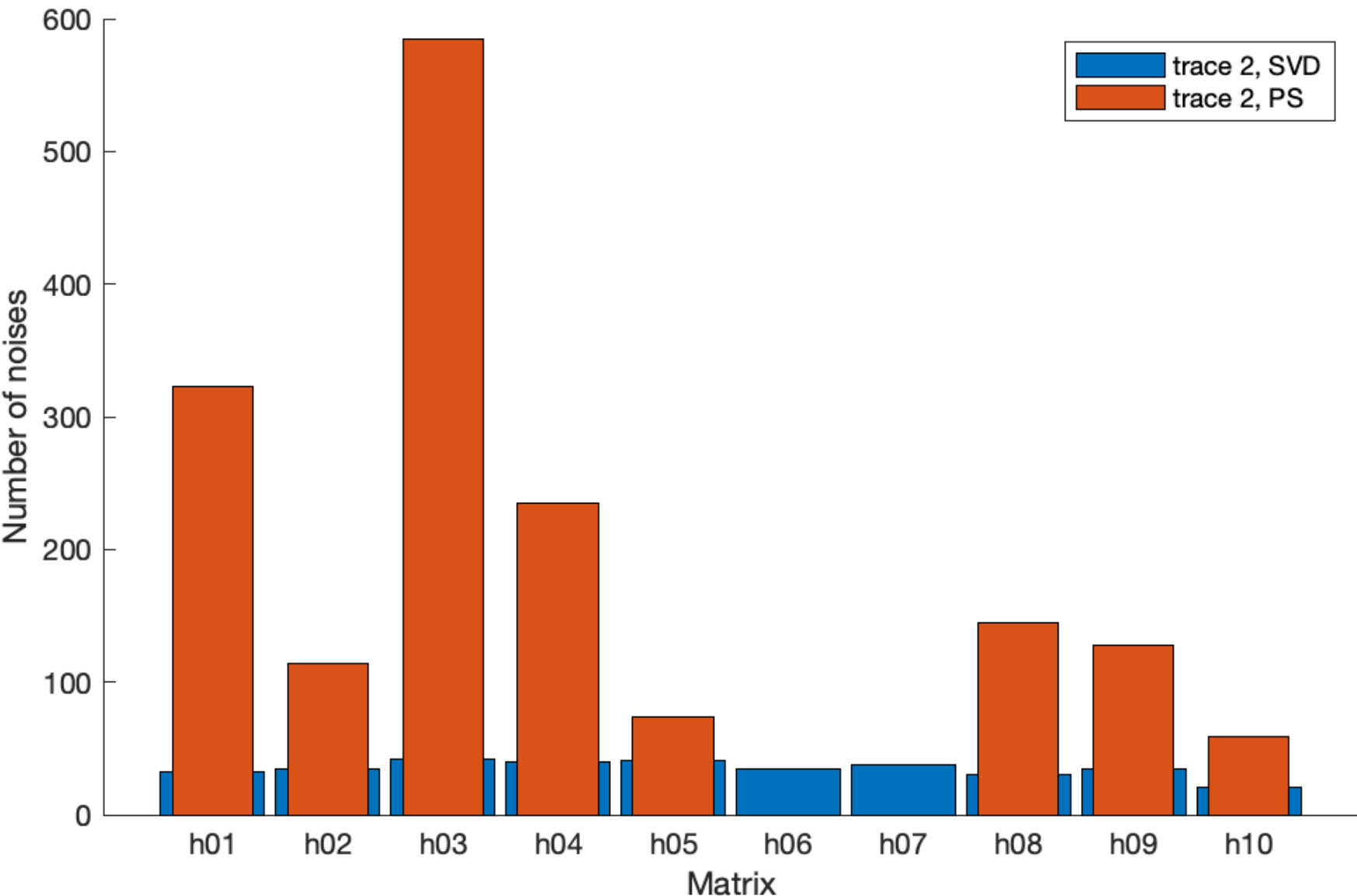
Noises required for each trace on 24^4 quenched lattice at κ_{crit} . Trace tolerance = 0.0005×24^4 .

Singular value deflation (SVD)

- Run the symmetric Lanczos algorithm on $\gamma_5 A$ (special for QCD) to generate singular values and vectors
- We test the full polynomial-subtracted, SVD trace,

$$Tr(A^{-1}) \approx \underbrace{Tr(A^{-1} - p(A))}_{(1)} + \underbrace{Tr\left(p(A) - \sum_i \frac{1}{\sigma_i} v_i u_i^\dagger\right)}_{(2)} + Tr\left(\sum_i \frac{1}{\sigma_i} v_i u_i^\dagger\right)$$

Matrix	nsvals
h01	62
h02	71
h03	69
h04	70
h05	73
h06	69
h07	61
h08	73
h09	56
h10	53

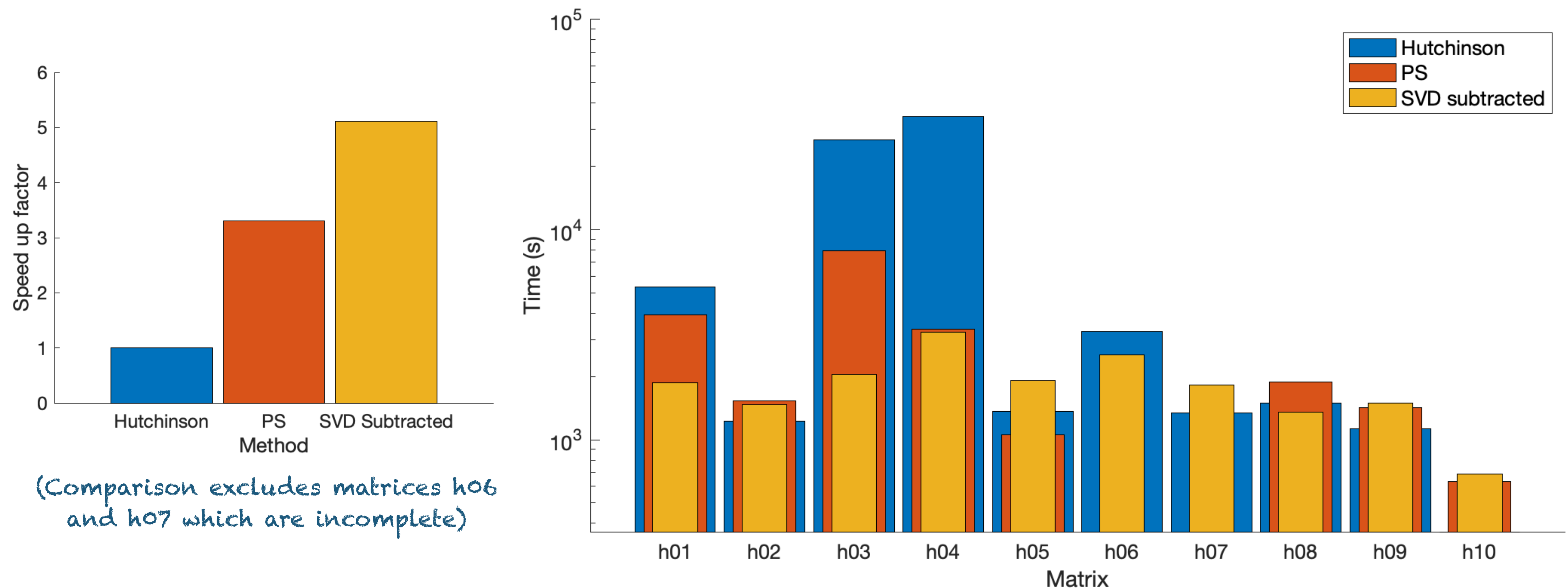


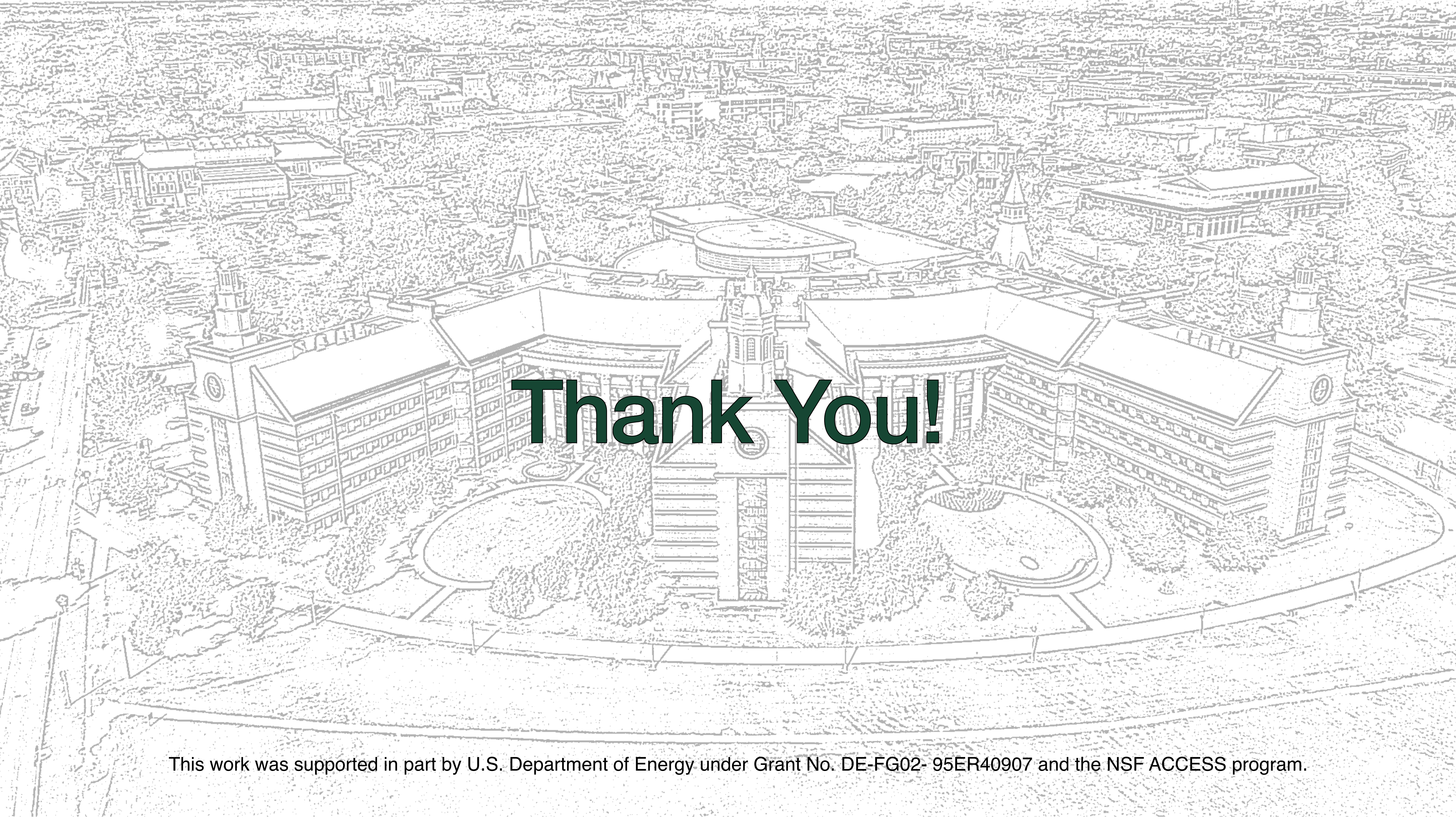
Noises required for each trace on quenched 24^4 lattice at κ_{crit} .
Trace tolerance = 0.0005×24^4 .

Trace times

Polynomial subtraction and SVD

- Compare the total times for regular Hutchinson's method, a polynomial subtraction trace, and a full polynomial subtraction and SVD subtracted trace.

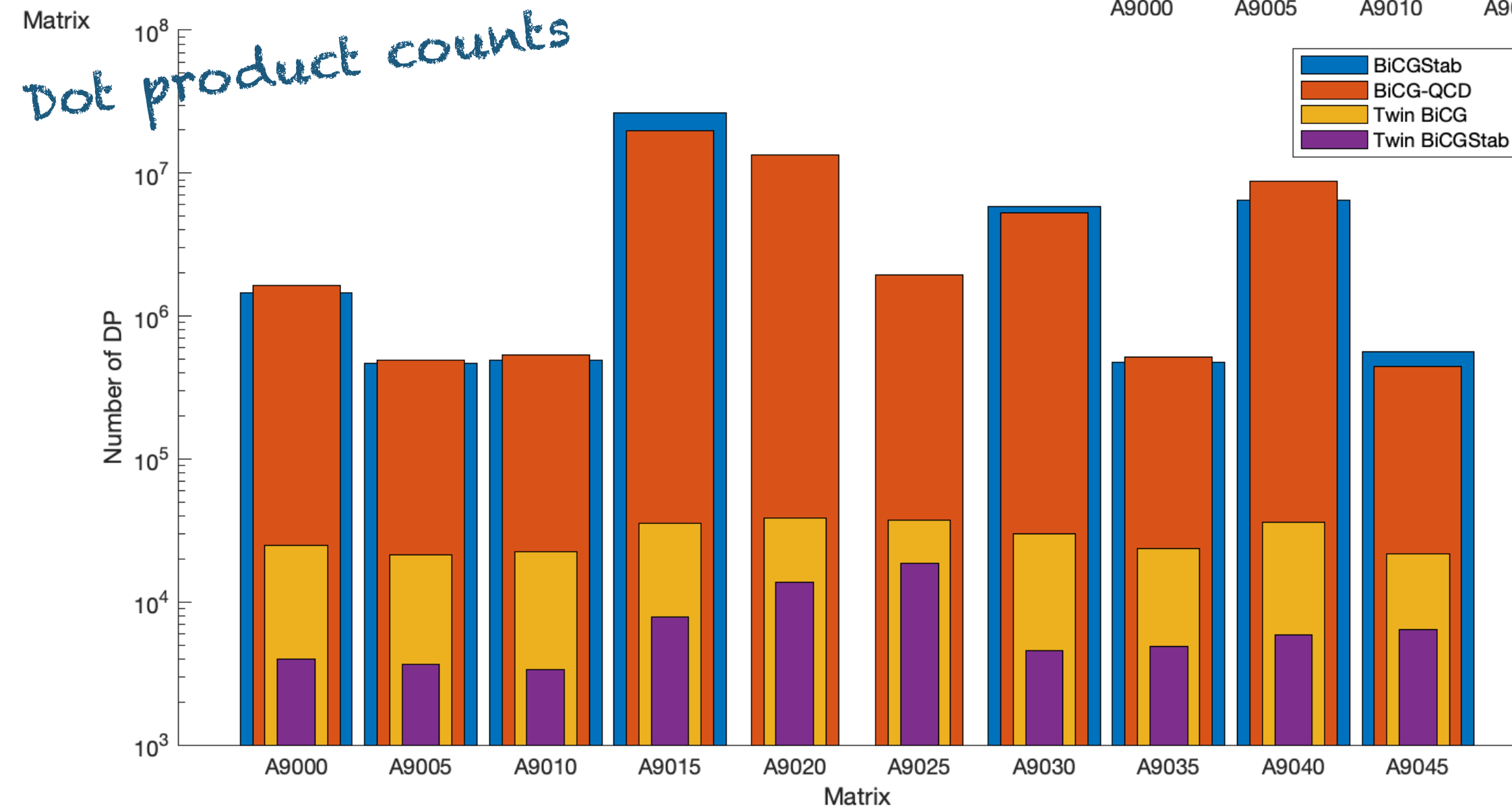
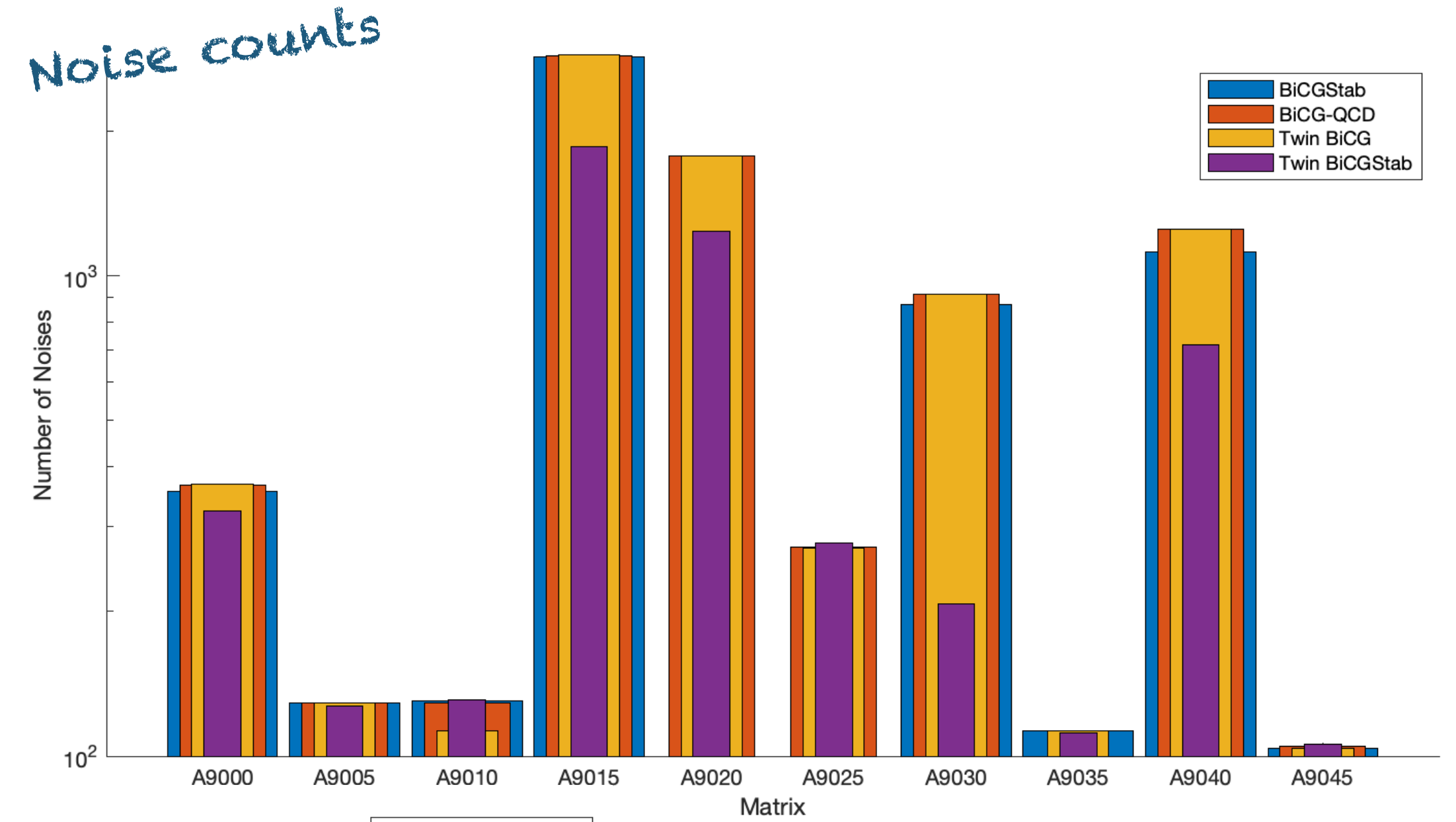
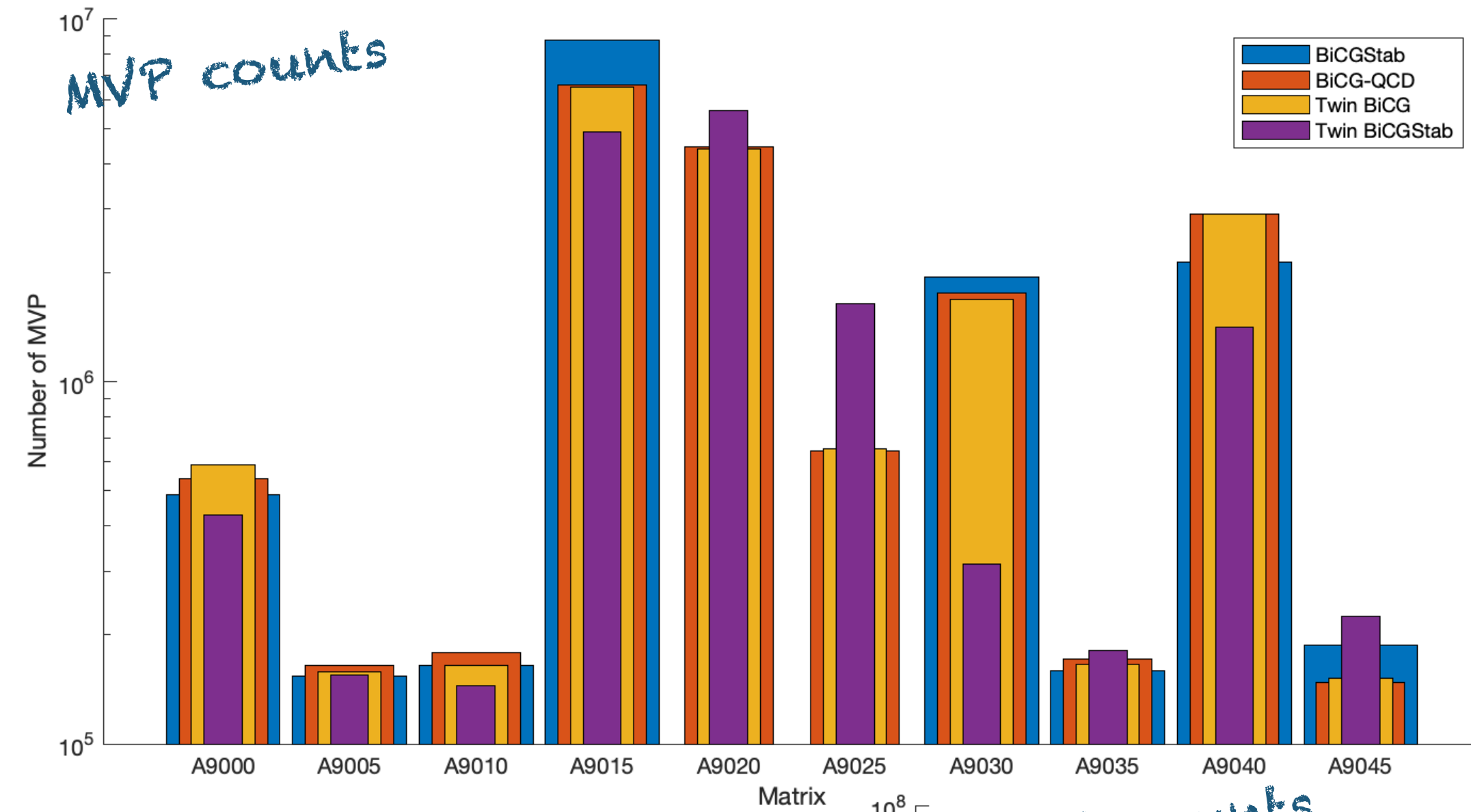




Thank You!

This work was supported in part by U.S. Department of Energy under Grant No. DE-FG02- 95ER40907 and the NSF ACCESS program.

A matrices, physical pion mass



B matrices, physical pion mass

